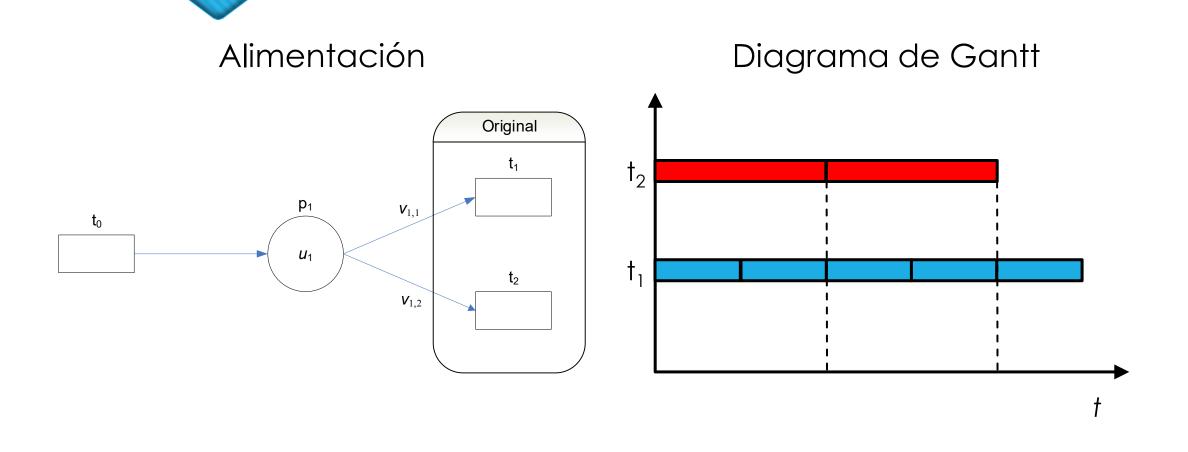


# Redes de Petri Parte III

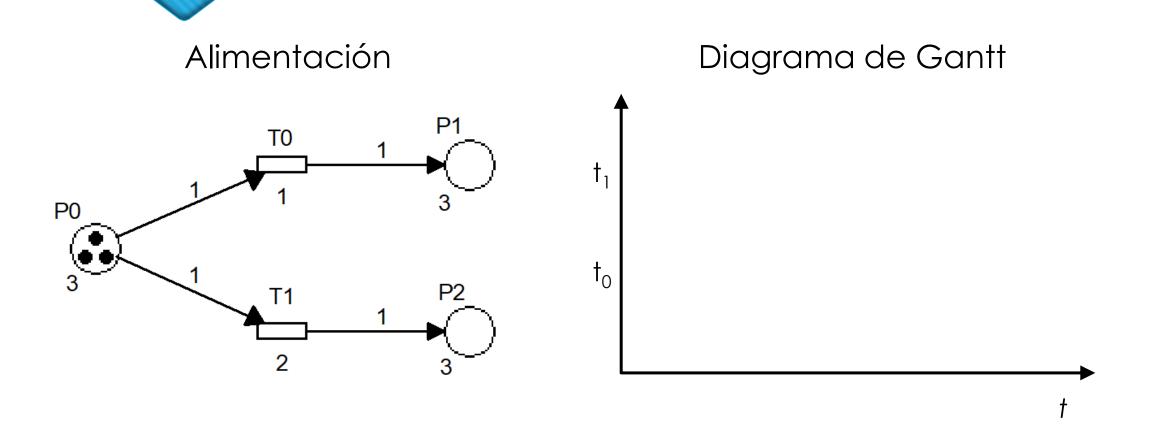
Enrique E. Tarifa, Facultad de Ingeniería, UNJu

# Estación de trabajo

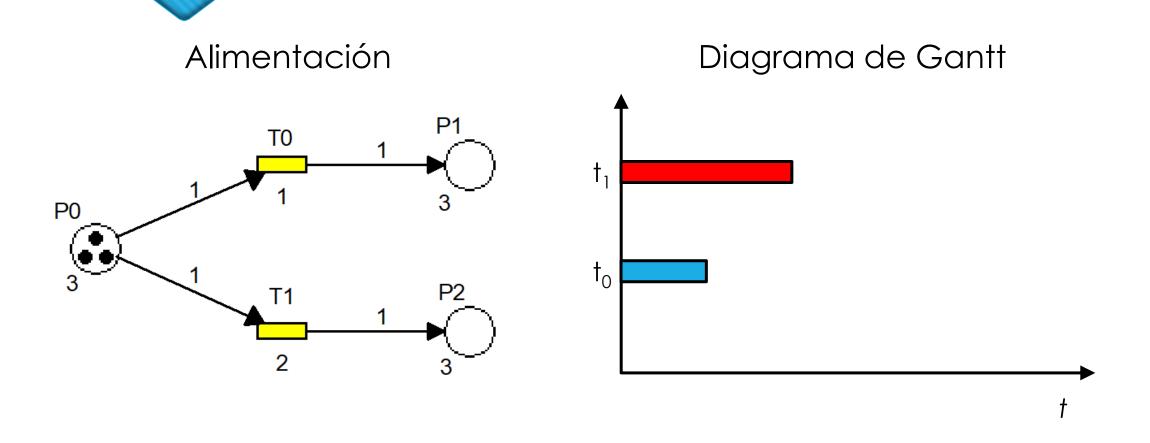


El lugar p<sub>1</sub> elige al azar.

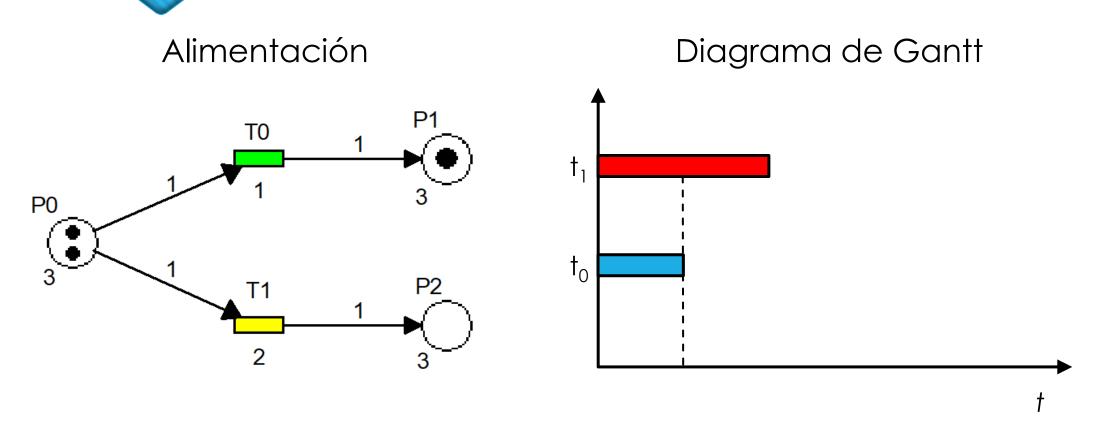
Conflicto.hps



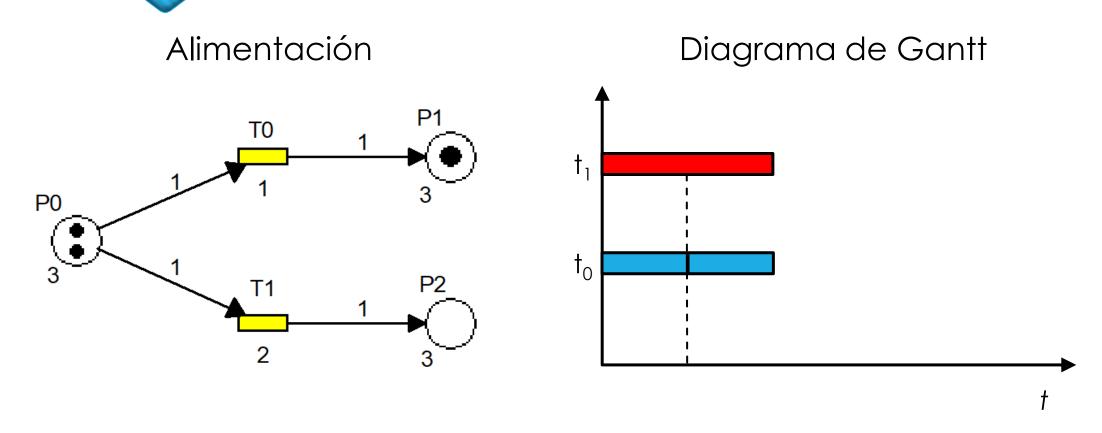
 $U_0 = 3$ ,  $U_1 = 0$ ,  $U_2 = 0$ , t = 0



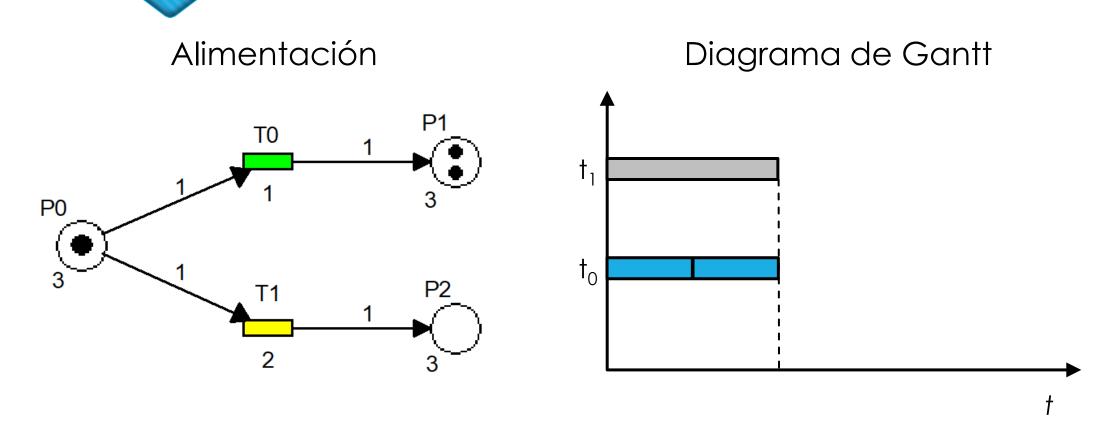
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$$U_0 = 2$$
,  $U_1 = 1$ ,  $U_2 = 0$ ,  $t = 1$ 



$$U_0 = 2$$
,  $U_1 = 1$ ,  $U_2 = 0$ ,  $t = 1$ 



$$U_0 = 1$$
,  $U_1 = 2$ ,  $U_2 = 0$ ,  $t = 2$ 

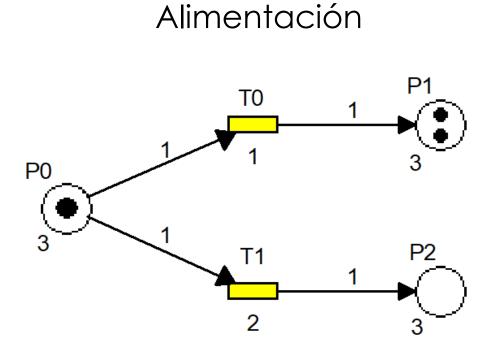
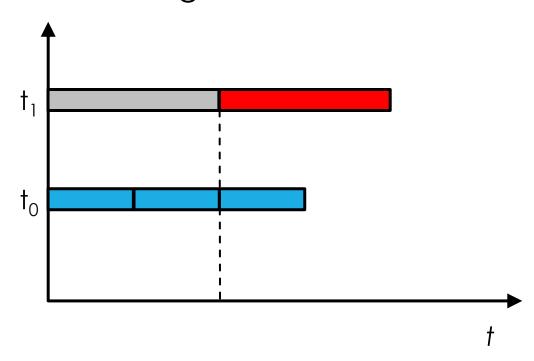


Diagrama de Gantt



$$U_0 = 1$$
,  $U_1 = 2$ ,  $U_2 = 0$ ,  $t = 2$ 

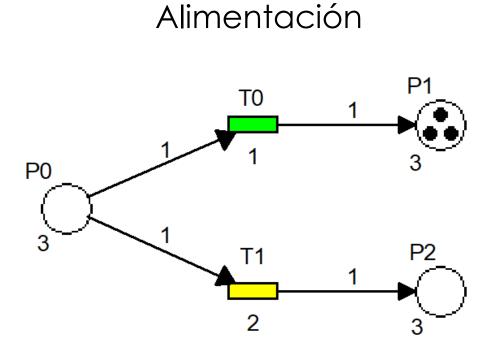
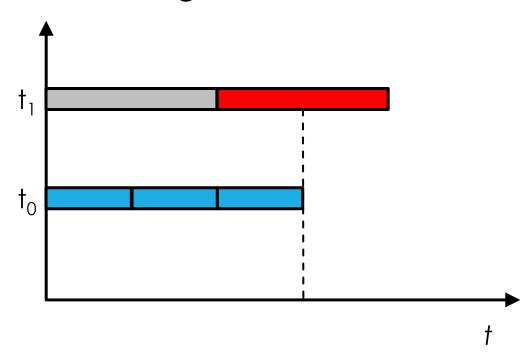


Diagrama de Gantt



$$U_0 = 0$$
,  $U_1 = 3$ ,  $U_2 = 0$ ,  $t = 3$ 

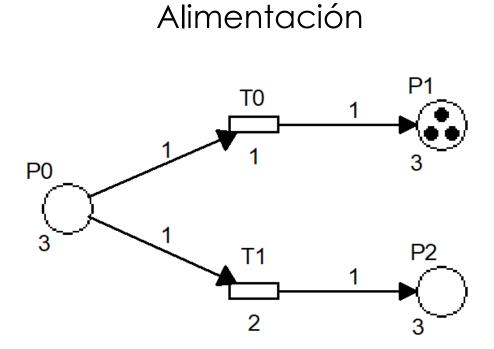
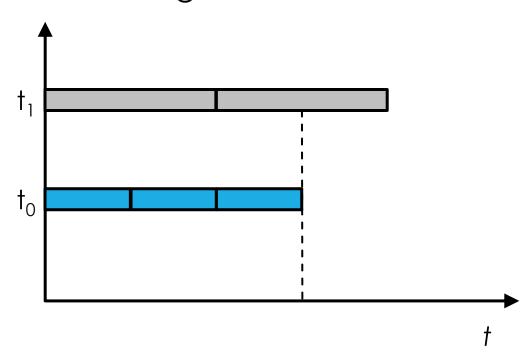
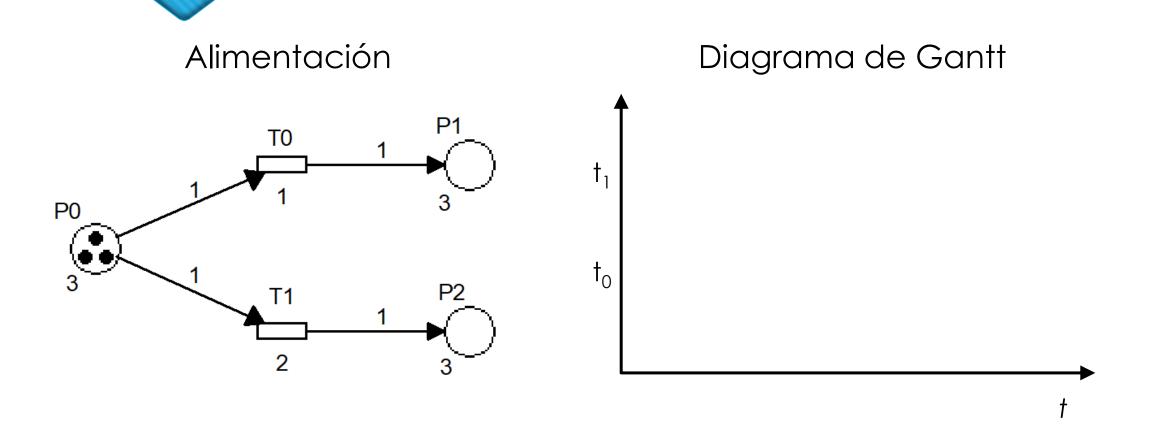


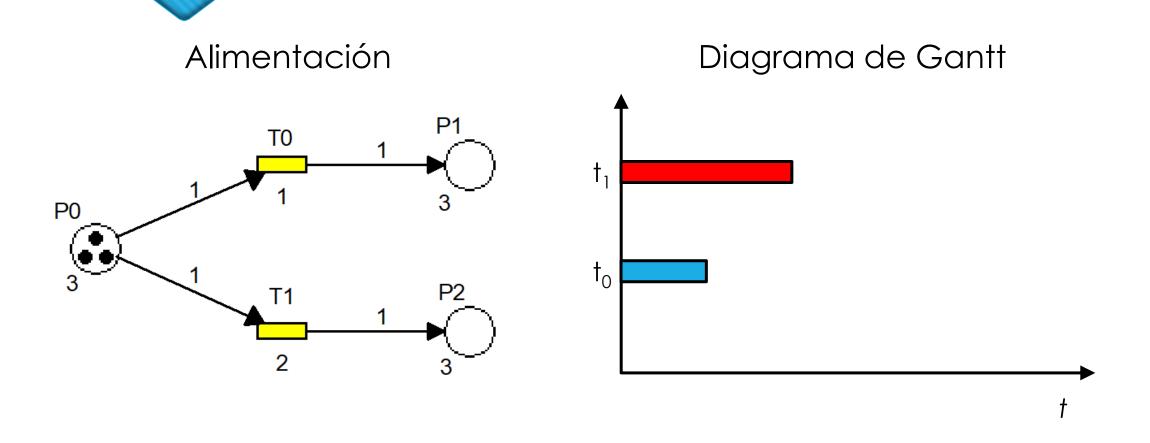
Diagrama de Gantt



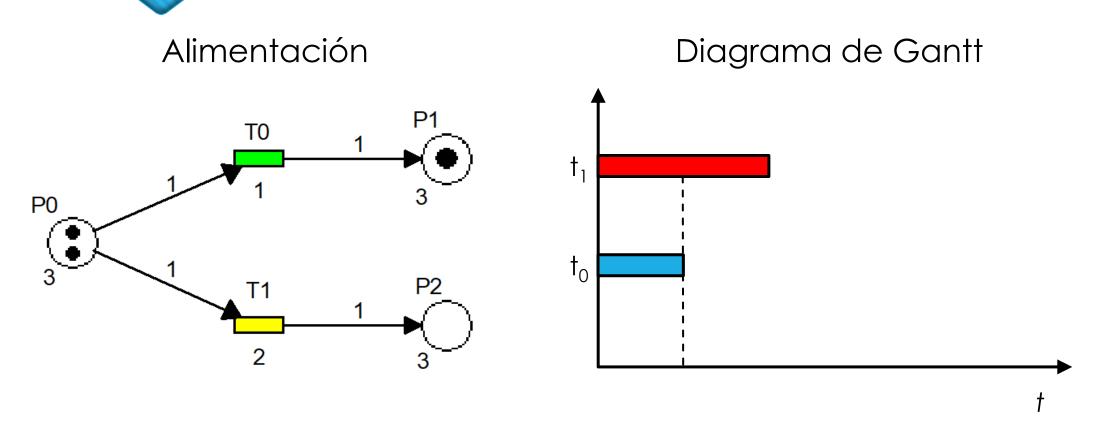
$$U_0 = 0$$
,  $U_1 = 3$ ,  $U_2 = 0$ ,  $t = 3$ 



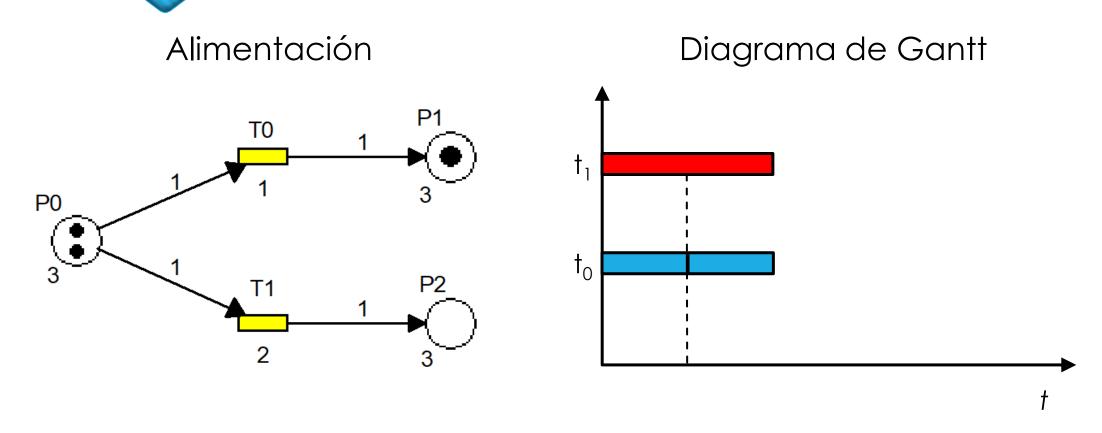
 $U_0 = 3$ ,  $U_1 = 0$ ,  $U_2 = 0$ , t = 0



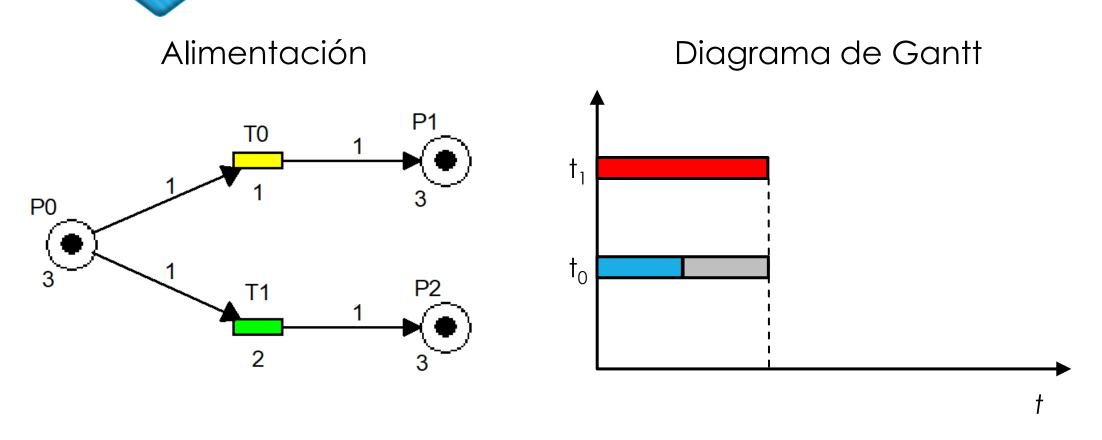
 $U_0 = 3$ ,  $U_1 = 0$ ,  $U_2 = 0$ , t = 0



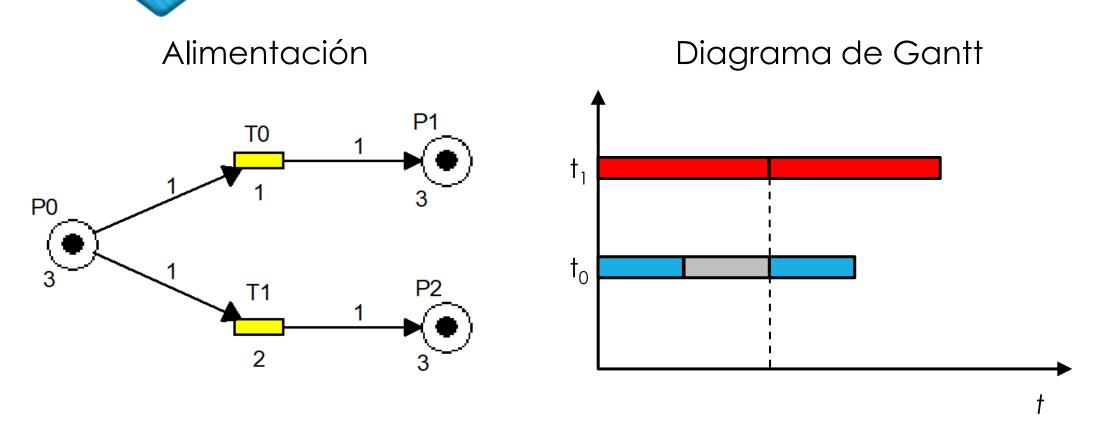
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,  $U_1 = 1$ ,  $U_2 = 0$ ,  $t = 1$ 



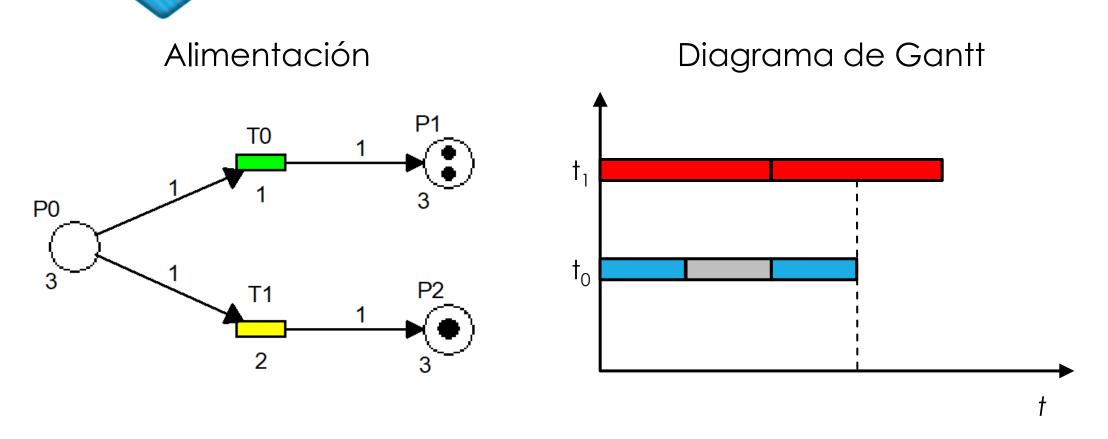
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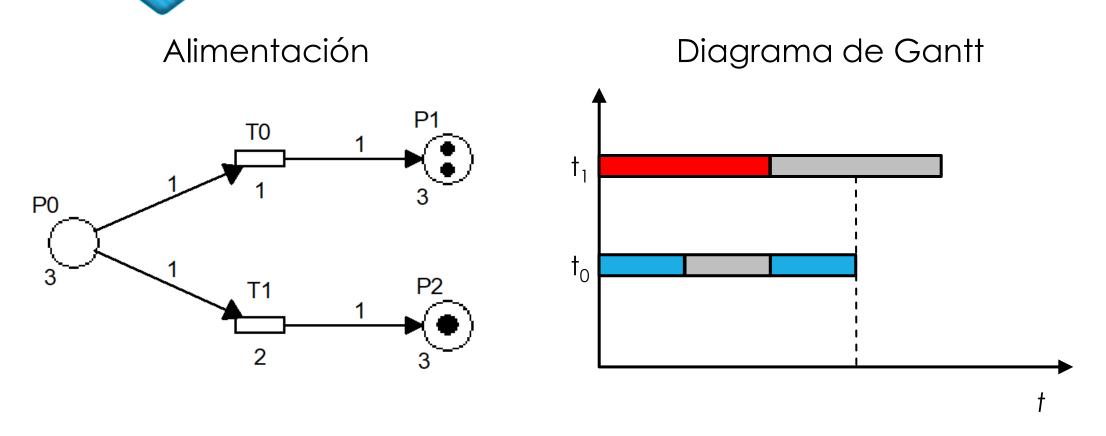
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,  $U_1 = 1$ ,  $U_2 = 1$ ,  $t = 2$ 



$$U_0 = 1$$
,  $U_1 = 1$ ,  $U_2 = 1$ ,  $t = 2$ 

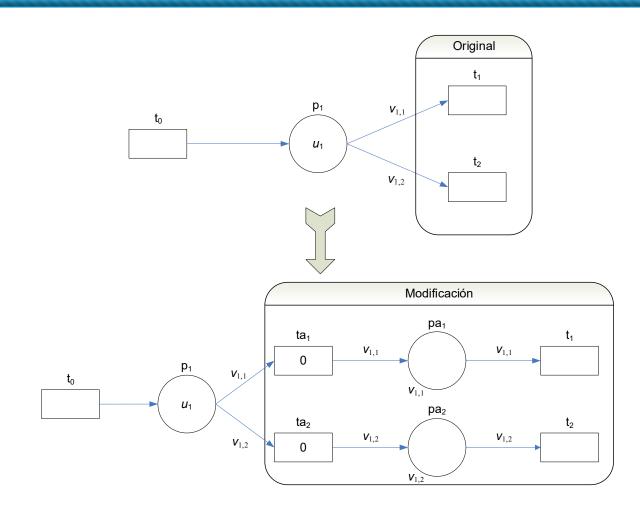


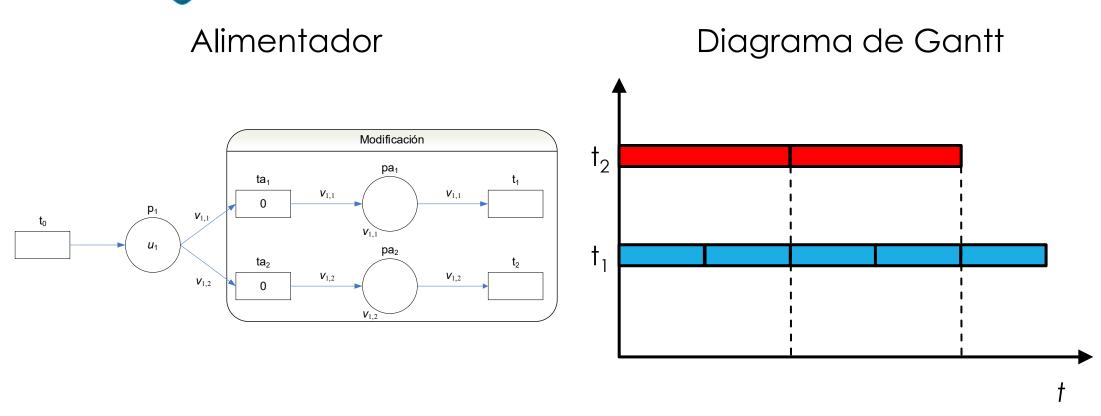
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,  $U_1 = 2$ ,  $U_2 = 1$ ,  $t = 3$ 



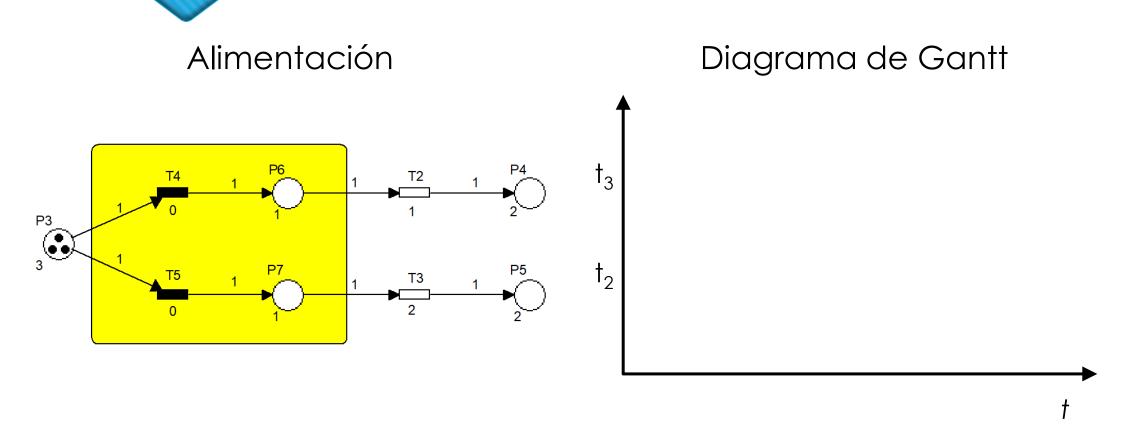
$$U_0 = 0$$
,  $U_1 = 2$ ,  $U_2 = 1$ ,  $t = 3$ 

#### Alimentador

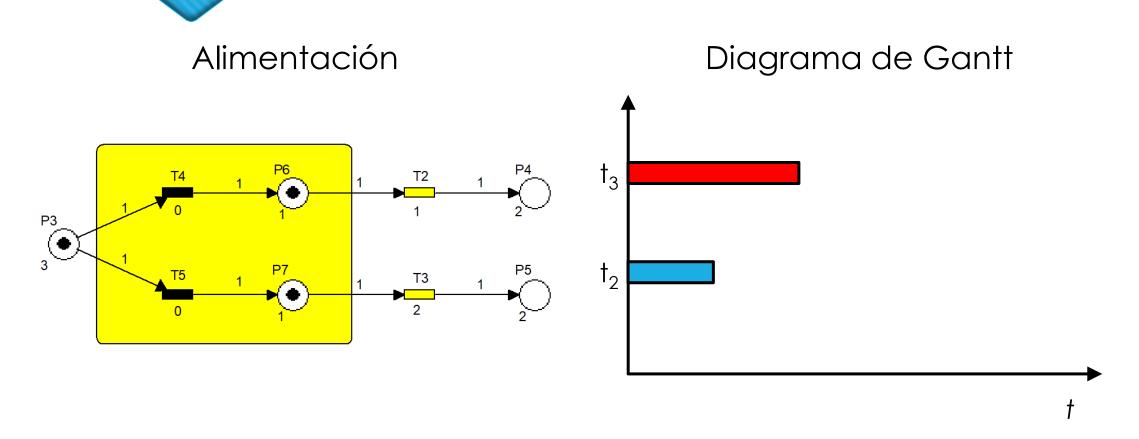




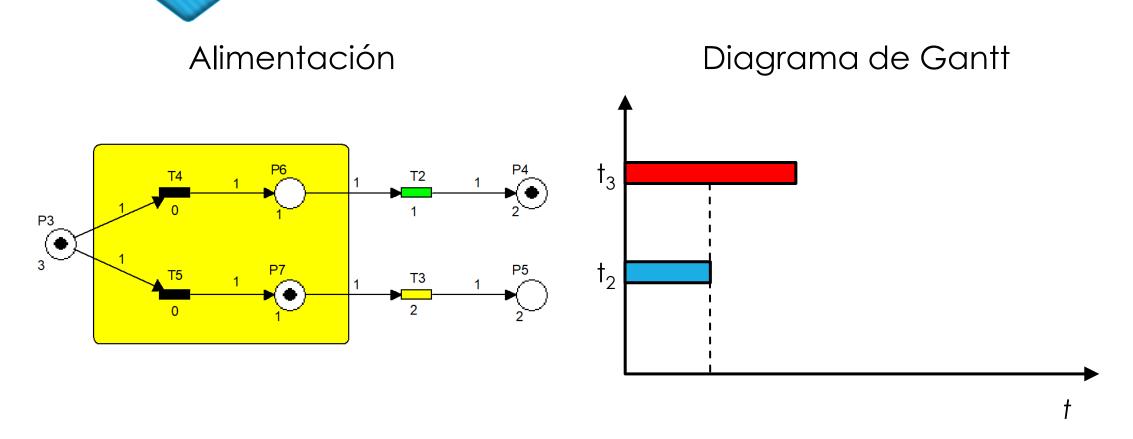
Si el lugar p<sub>1</sub> no tiene suficientes marcadores, elige al azar.



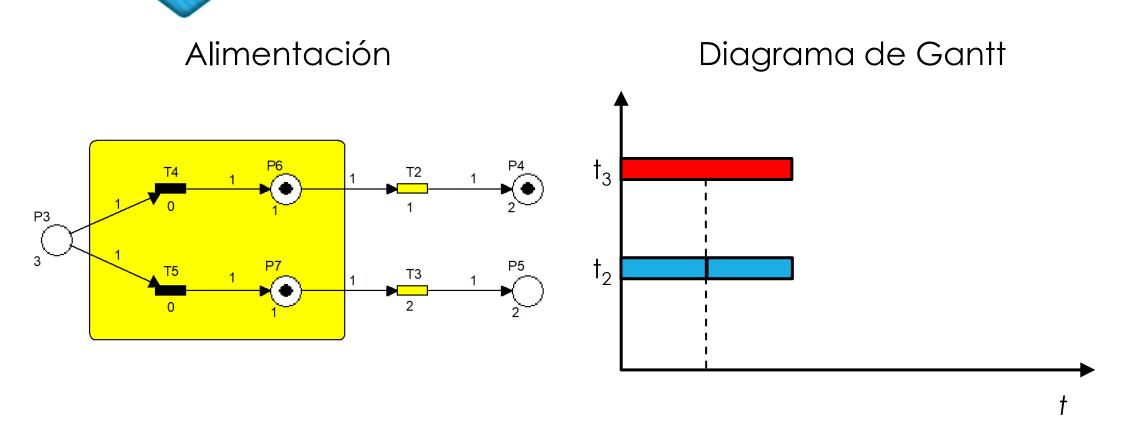
$$U_3 = 3$$
,  $U_4 = 0$ ,  $U_5 = 0$ ,  $t = 0$ 



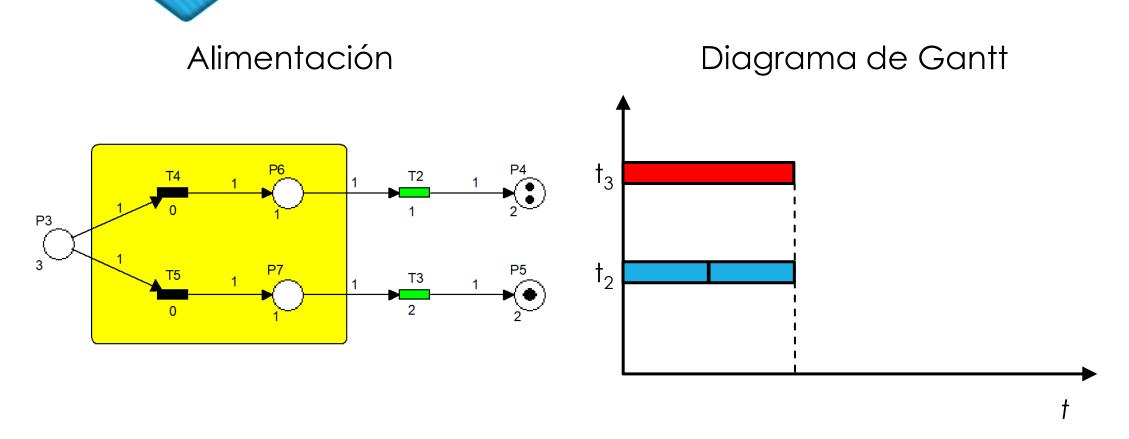
$$U_3 = 1$$
,  $U_4 = 0$ ,  $U_5 = 0$ ,  $t = 0$ 



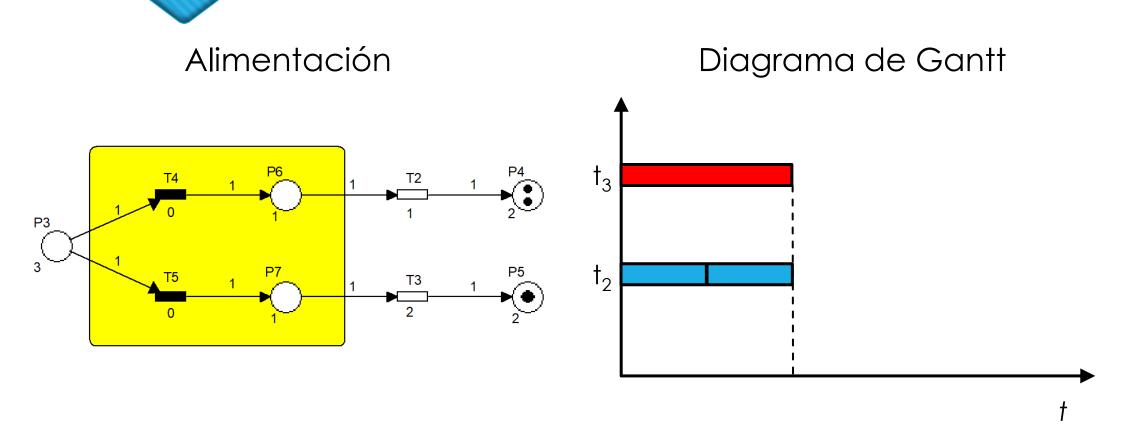
$$U_3 = 1$$
,  $U_4 = 1$ ,  $U_5 = 0$ ,  $t = 1$ 



$$U_3 = 0$$
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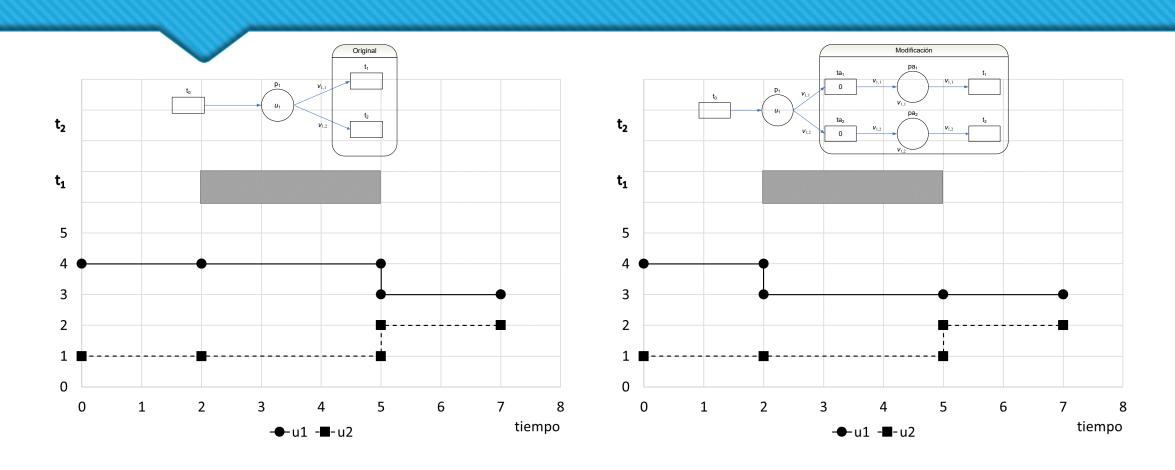


$$U_3 = 0$$
,  $U_4 = 2$ ,  $U_5 = 1$ ,  $t = 2$ 

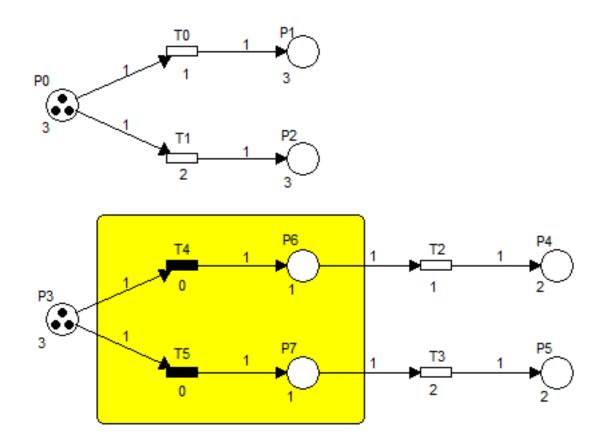


$$U_3 = 0$$
,  $U_4 = 2$ ,  $U_5 = 1$ ,  $t = 2$ 

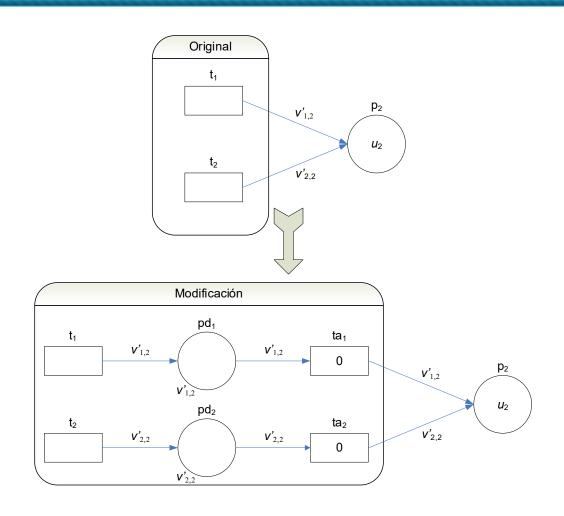
## Diagrama de Gantt y evolución



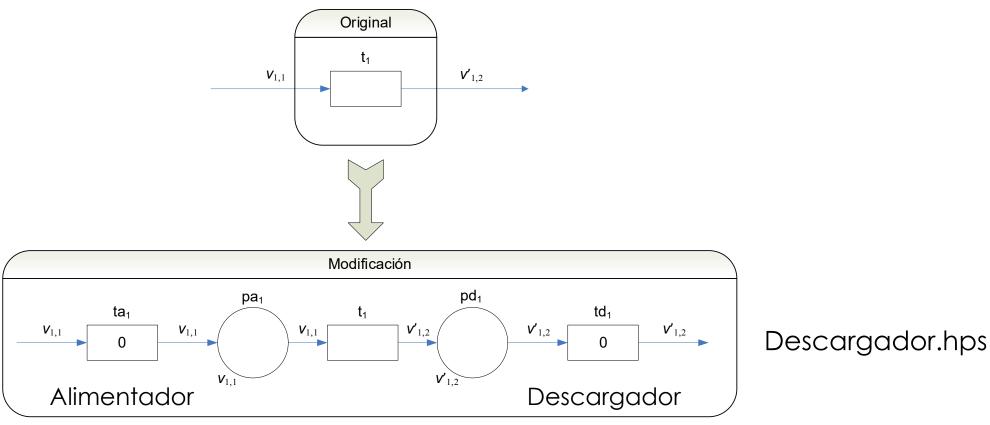
#### Alimentador



# Descargador



# Estación de trabajo

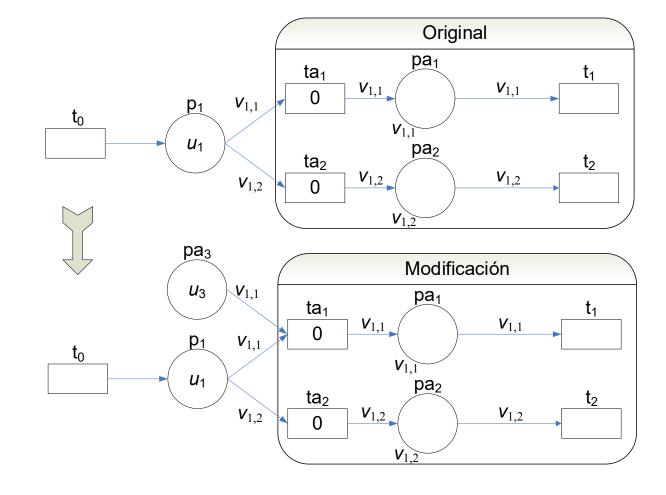


La estación de trabajo puede acortar el tiempo de producción.

# Orden de producción

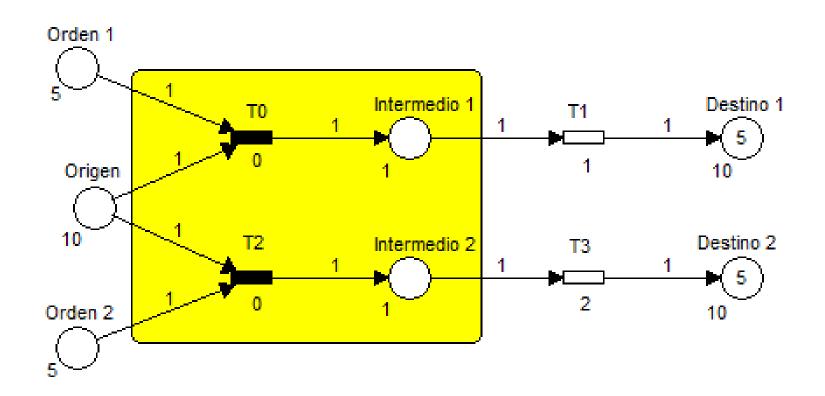
#### Orden de producción

- $ot_1 < t_2$
- υ<sub>3</sub>: Cantidad de marcadores destinados a la línea más rápida.
- Hay que colocar algo similar en la línea lenta.



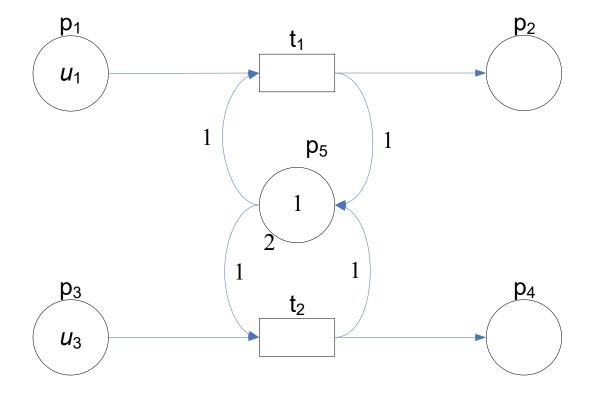
Orden de producción.hps

# Orden de producción

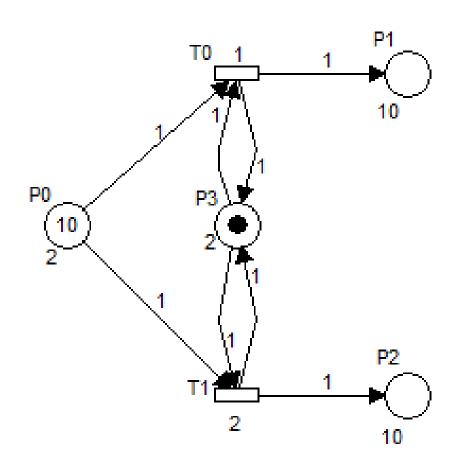


# Asignación de recursos

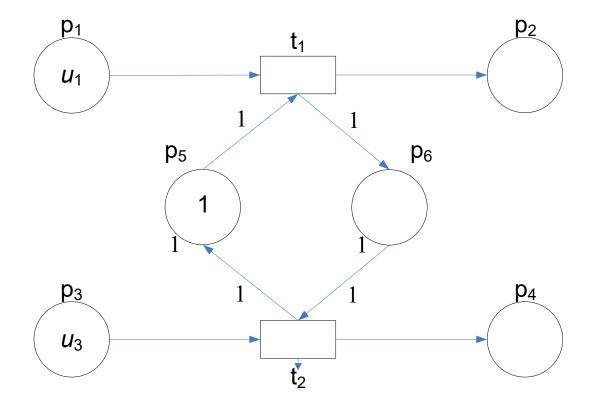
# Operario al azar



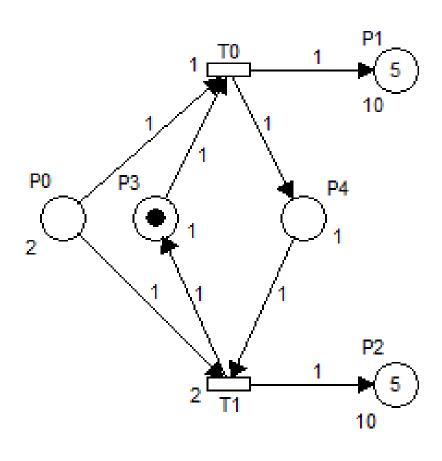
# Operario al azar



# Operario alternativo



# Operario alternativo



#### Determinación de recursos

- $\circ$  El tiempo original de la actividad es  $t_1$ .
- Los recursos se multiplican por n.
- Si la tarea se puede dividir,  $t_n = t_1/n$ .

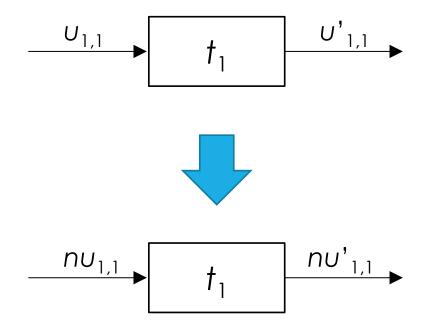






#### Determinación de recursos

- $\circ$  El tiempo original de la actividad es  $t_1$ .
- Los recursos se multiplican por n.
- Si la tarea no se puede dividir,
  multiplicar por n las multiplicidades.

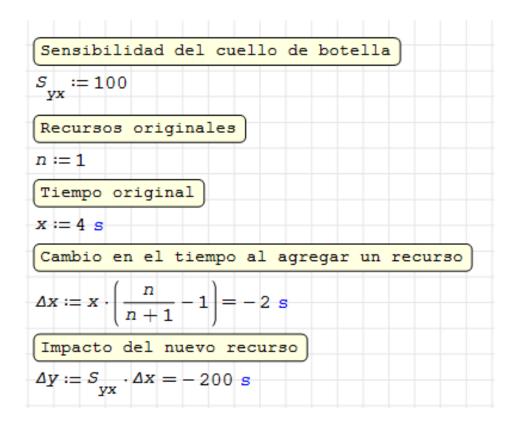


#### Asignación de recursos

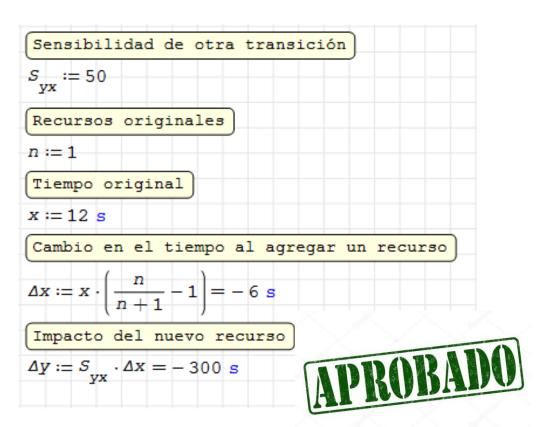
- Un nuevo recurso se debe asignar a la transición donde tenga el máximo impacto.
- Impacto en  $\Delta y = S_{y,x} \Delta x$
- O No siempre el máximo impacto estará en el cuello de botella:
  - O Cuello de botella:  $S_{y,x} = 100$ ,  $\Delta x = 2$  s,  $\Delta y = 200$  s
  - Otra transición:  $S_{y,x} = 50$ ,  $\Delta x = 6$  s,  $\Delta y = 300$  s

### Asignación de recursos

#### Cuello de botella



#### Otra transición



# Aplicaciones

#### Aplicaciones

- o Determinación de la producción
- Determinación del tiempo de producción
- O Determinación del estado en un tiempo dado
- Dimensionamiento de depósitos
- Identificación de inventarios críticos
- O Determinación del inventario mínimo
- Identificación de cuellos de botella
- Evaluación de asignación de recursos

